Recent algorithm and machine developments for lattice QCD

Ken-Ichi Ishikawa (Hiroshima Univ.) Lattice 2008, July 16



1. Plan of My Talk

2. Machine trends

- New machines
 - BG/P, T2K, QPACE project, Pet-Ape project.
- Many cores
 - GPGPU CUDA
- 3. Algorithmic developments for dynamical QCD (Wilson type) O HMC with
 - Preconditioning for HMC action and UV/IR separation: Domain-Decomposition,RHMC, Schur complement...
 - Multiple timescale MD integrator
 - Solver with
 - Inner-Outer(mixed prec.), Deflation, Adoptive Multi Grid.
- 4. Outlook: Physics at 1Pflops
 - Finer lattice (continuum limit or charm quark)
 - Larger volume (multi hadron system)

- 2. Machine Trends
 - New machines
 - ○Blue Gene/P



Successor of QCDSP, QCDOC, Blue Gene/L

[P.Boyle et al., IBM J. Res. and Dev. 49 (2005)

http://www.research.ibm.com/journal/rd/492/boyle.html]

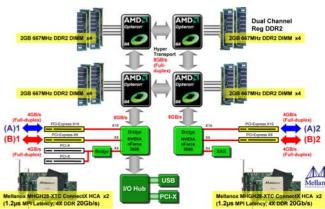
- 4Way SMP PowePC@0.85GHz
- Scalable 3D torus network
- Population is incleasing
- Thin node / O(100,000) Many nodes
- Byte/Flop balanced
- Fine grained parallelization.

- 2. Machine Trends (cont'd)
- New machines
 - T2K open super computer project (Tsukuba-Tokyo-Kyoto)



[http://www.open-supercomputer.org/]

- 4 Way Opteron (Barcelona) node cluster (commodity base).
- 648nodes@tsukuba, 147GFlops/node (Fat node)
- Quad core, 4 way
- Multi-rail fat tree network
- Many core / Fat node / O(1,000) few nodes
- Maintain Byte/Flop at each level
- Data Blocking is required



New machines (for QCD)

QPACE project (QCD PArallel computing on the CEII/B.E.)
 2008-2009 [Poster by A. Nobile "Status of the QPACE Project"]

- Fund by Deuche Forschungsgemeinscaft (DFG)
- Collaboration with IBM Germany.
- Dedicated for LQCD. 200TFlops (2009)
- Cell Broadband Engine cluster. [PowerXCell 8i, 102GFlops(DP)]
- Custom 3D torus Scalable network (FPGA)
- Low power consumption 1.5W/GFlops
- Many core / Fat node / O(1,000) Few nodes
- Maintain Byte/Flop at each level
- Data Blocking is required

QCD with CELL: Spary,Hill,Trew hep-lat/0804.3654; S.Motoki & A. Nakamura Lat2007; F.Belletti et al. LAT2007



- 2. Machine Trends (cont'd)
- New machines

Pet-APE project (Petaflops Array Processor Experiment)

[INFN APE Groupe, Italy, to apper in NUOVO CHIMENTO]

EMAIL From Davide Rossetti@ROMA1.INFN

- Successor of APEmille, apeNEXT.
- Reference computing platform for LQCD (2009-2014)
- Custom CPU: Apotto
- Custom Network ApeNet+ 3D Torus.
- Aiming for good price/performance.
- Thin node / O(100,000) Many nodes ?

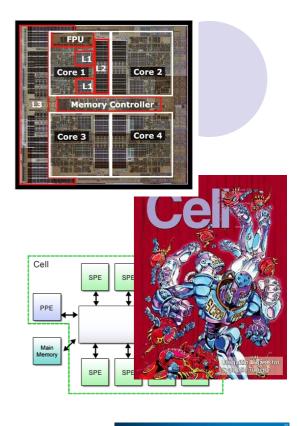
Many cores (QPACE, T2K)

To make use of the full machinery of many transistors on a chip, many core architecture is employed for recent processor

- Intel: Core 2 Quad (4cores, 3GHz, 48GFlops),...
- AMD: Phenom (4cores, 2.4GHz, 38GFlops), ...
- IBM: Power X Cell 8i (1+8cores, 3.2GHz, 102GFlops)
- SUN: UltraSparc T2 (8cores)

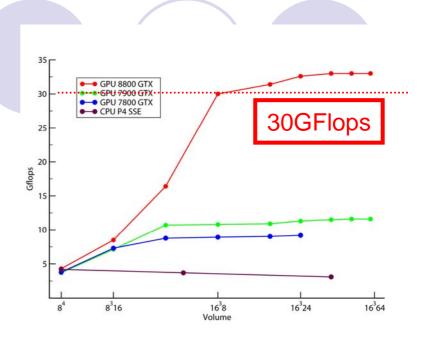
The trend is 8 cores, 16 cores,...., many cores

- Intel larrabee 80 cores?
- AMD/ATI GPGPU firestrem 800 cores?
- NVIDIA GPGPU CUDA 240 cores?
- As a many core example , GPGPU





 "Lattice QCD as a video game",
 G.I.Egri, Z.Fodor, S.D.Katz, D.Nogradi, K.K.Szabo, hep-lat/0611022.
 NVIDIA G80 arch. > 300 GFlops(SP)
 Lattice Wilson kernel > 30 GFlops
 Difficult to program using Graphic API (OpenGL)



- NVIDIA provides HPC GPGPU language
 - CUDA (a C/C++ simple extension)
 - Easy to learn, but requires hardware/memory model knowledge

[Poster by C. Rebbi, "Blastign Through Lattice Calc. using CUDA" talk by F. Di Renzo, "GPU computing for 2-d spin systems:CUDA vs OpenGL"]

 My experience with CUDA (GeForce 8800 GTX)
 [NO WARRANTY CUDA code: http://theo.phys.sci.hiroshimau.ac.jp/~ishikawa/CUDA/CudaQCDSolver_0.06.tar.gz]

Hopping matrix mult (16^4) can also achieve > 40 GFlops.

C. Rebbi (Poster): Wilson Dirac 62GFlops! with Nvidia Tesla C870

My experience with CUDA (GeForce 8800 GTX)

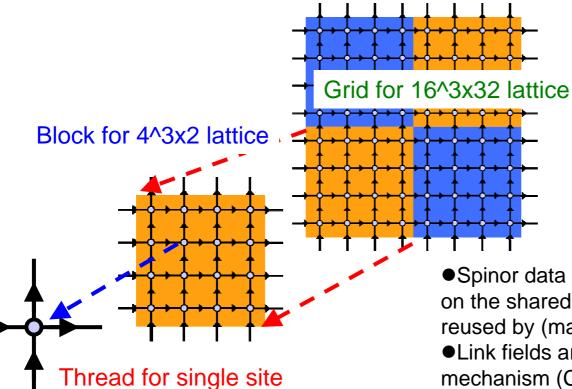
```
static __inline__ _device__
void matvec2(float *ur, float *ui, float *yr, float *yi, float
*uyr, float *uyi){
 float ux3r,ux3i;
// 0-2
 ux3r = ur[1+COL*0]*ur[2+COL*1] + ui[2+COL*0]*ui[1+COL*1];
 ux3i = ur[2+COL*0]*ui[1+COL*1] + ui[2+COL*0]*ur[1+COL*1];
 ux3r = -ux3r;
 ux3i = -ux3i;
 ux3r += ui[1+COL*0]*ui[2+COL*1] + ur[2+COL*0]*ur[1+COL*1];
 ux3i += ur[1+COL*0]*ui[2+COL*1] + ui[1+COL*0]*ur[2+COL*1];
 ux3r = -ux3r;
 ux3i = -ux3i;
 *(uyr+0)
            = *(ur+0+COL*0) * *(yr+0);
  *(uyr+0)
           += *(ur+0+COL*1) * *(yr+1);
  *(uyr+0)
           += ux3r
                            * *(yr+2);
 *(uyr+0)
           = - *(uyr+0);
  *(uvr+0)
           += *(ui+0+COL*0) * *(vi+0);
           += *(ui+0+COL*1) * *(yi+1);
  *(uyr+0)
 *(uyr+0)
                            * *(yi+2);
           += ux3i
 *(uvr+0)
           = - *(uvr+0);
 *(uyr+0+COL) = *(ur+0+COL*0) * *(yr+0+COL);
 *(uvr+0+COL)+= *(ur+0+COL*1) * *(vr+1+COL);
 *(uyr+0+COL)+= ux3r
                           * *(yr+2+COL);
  *(uyr+0+COL) = - *(uyr+0+COL);
  *(uvr+0+COL)+= *(ui+0+COL*0) * *(vi+0+COL);
  *(uyr+0+COL)+= *(ui+0+COL*1) * *(yi+1+COL);
  *(uyr+0+COL)+= ux3i
                     * *(yi+2+COL);
  *(uvr+0+COL) = - *(uvr+0+COL);
  *(uyi+0) = *(ur+0+COL*0) * *(yi+0);
  *(uyi+0) += *(ur+0+COL*1) * *(yi+1);
  *(uyi+0) += ux3r * *(yi+2);
  *(uyi+0) += *(ui+0+COL*0) * *(yr+0);
  *(uyi+0)
           += *(ui+0+COL*1) * *(yr+1);
 *(uvi+0)
           += ux3i
                             * *(yr+2);
  *(uyi+0+COL) = *(ur+0+COL*0) * *(yi+0+COL);
 *(uyi+0+COL)+= *(ur+0+COL*1) * *(yi+1+COL);
  *(uvi+0+COL)+= ux3r
                           * *(vi+2+COL);
  *(uyi+0+COL)+= *(ui+0+COL*0) * *(yr+0+COL);
```

Cuda code example: Link variable times 2-Spinor code almost C language

For Single site data, w, y: 2 - spinor, U: SU(3) matrixfor $\alpha = 1,2$ and a = 1,2,3 $w(a, \alpha) = \sum_{b=1}^{3} U(a,b) y(b, \alpha)$

My experience with CUDA (GeForce 8800 GTX)

CUDA Programming model
Single Program Multiple Data (SPMD)
Nested threading. Grid / Block / Thread
Thread ID + Block ID (Corresponds to MPI RANK)
Block has local memory shared by threads in a block.



се					
	Grid 1				
	Block	Block	Block		
	(0, 0)	(1, 0)	(2, 0)		
	Block	Block	Block		
	(0, 1)	(1, 1)	(2, 1)		

Block (1, 1)			
Thread (0, 0)	Thread (1, 0)	Thread (2, 0)	Thread (3, 0)	Thread (4, 0)
Thread (0, 1)	Thread (1, 1)	Thread (2, 1)	Thread (3, 1)	Thread (4, 1)
Thread (0, 2)	Thread (1, 2)	Thread (2, 2)	Thread (3, 2)	Thread (4, 2)

•Spinor data are vector loaded [100GByte/sec] on the shared memory on each block. They are reused by (max 8 times/ min 4 times).

•Link fields are loaded via Texture Fetching mechanism (Cached).

Dev

My experience with CUDA (GeForce 8800 GTX)

CUDA Language is Ready for Lattice QCD!!!

•CELL has similar feature?

•How about other accelerator? (AMD/ATI card, ClearSpeed)

CUDA sample solver residual history [D.D.-preconditioned BiCGStab] (Lattice size: $16^3 \times 32$, block size: $4^3 \times 2$) 10⁰ 10-1 10⁻² Residual |b-Dx|/|b| 10⁻³ 10⁻⁴ 10⁻⁵ 10⁻⁶ 10-7 10⁻⁸ 0 5 10 15 20 25 30 35 40 45 50 iter.

A test result with CUDA solver (Single precision)

See also CUDA Works in This conference:

•F. Di Renzo, "GPU computing for 2-d spin sytems:CUDA vs OpenGL"

•C. Rebbi, "Blasting Through Lattice Calculations using CUDA"

CELL Works:

•V. Kindratenko, "Cell processor implementation of a MILC lattice QCD application"

This year Nvidia and AMD/ATI provide DP enabled architecture

- NVIDIA GT200 (Tesla 10series)
 - 240 SP (SP cores), 30 DP cores
 - ~1,000(or 600)Glops(SP), ~90GFlops(DP)

C. Rebbi (Poster): Wilson Dirac 100 GFlops! with Nvidia GTX280



- AMD/ATI RV770 (Firestream 9250)
 - 640 SP units, (160 DP units?)
 - 1.2TFlops (SP), 200 GFlops (DP)
 - AMD Stream SDK



Product: AMD FireStream[™] 9250 (Available September 2008)

Breaking the 1 TFLOPS barrier at under 150 watts!

○ For QCD

- No ECC, check the result on the host side.
- O(1000) thread programming/SIMD programming is required. (1site=1thread)
- Make use of the Local memories attached each core for good efficiency.
- Host ⇔device communication is limited by PCI-E x16 G2 speed (8GB/sec (sustained at 2GB/sec))

- Thin node / O(100,000) nodes (BG/P, Pet-APE)
 Uniform Fine Grained Parallelization is required.
 10GFlops/CPU, 100,000 nodes = 1PF
- Many cores / Fat node / O(1,000) nodes (T2K, QPACE, GPGPU)
 - Core/CPU/Node Hierarchy exists.
 - Data Bandwidth is not uniform.
 - Data blocking is required at each level.
- 1~10 PFlops machine trends?
 - My expectation is Many core/Fat node/O(1,000-10,000) nodes
 - O Near future: Intel larrabee, CELL, GPGPU,
 - 200 GFlops/CPU, 8 CPU/node=1.6TF/node, 1,000 node=1.6PFlops
 - 1 TFlops/CPU, 4 CPU/node = 4TF/node, 1,000 node= 4PF

- 3. Algorithmic developments for dynamical QCD
 - Recent improvement strategy (HMC)
 - Two key technologies for HMC algorithm
 - (1) Transform/split det[D] using preconditioner (Action Prec.)

UV/IR separation [de Forcrand, Takaishi, NPB(Proc.Suppl.)53,Lat96]

D: Lattice Dirac op., *P*: a preconditioner choose *P* s.t. cond(DP) < cond(D) and easy to compute det[P] $det[D] = \frac{det[DP]}{det[P]}$

Reduction of condition number of *D*

remove/suppress UV modes of D

DP: Preconditioned op. IR part/IR physics

P: Preconditioner UV part/UV physics

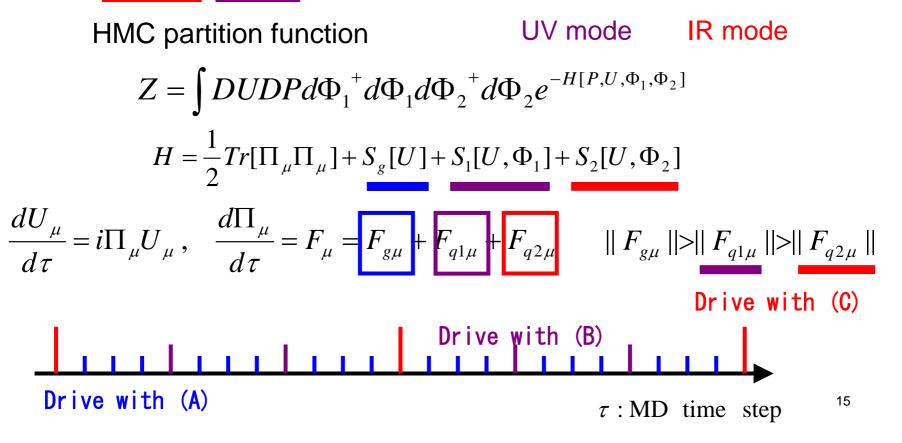
$$Z = \int \prod dU \det[D] e^{-S[U]} = \int \prod dU \det[DP] \det[P^{-1}] e^{-S[U]}$$

3. Algorithmic developments...(cont'd)

Recent improvement strategy (HMC)
 (2) Multi time step MD integrator

Multi time step MD integ. [Sexton-Weingarten, NPB 380(92)]

$$\det[DP]\det[P^{-1}] = \int d\Phi_1^{+} d\Phi_1 d\Phi_2^{+} d\Phi_2 e^{-\Phi_1^{+}P\Phi_1 - \Phi_2^{+}(DP)^{-1}\Phi_2}$$



- 3. Algorithmic developments...(cont'd)
- (1) Transform/split det[D] using preconditioner (Action Prec.)(A) Hasenbusch's heavy mass preconditioner

det[D] = det[D/D']det[D']D' has havy mass than D, and $D/D' \approx 1$ IR mode UV mode (B) Geometric preconditioner (Domain Decomposition) [Lüscher, JHEP 0305 '03, CPC 165 '05] $\det[D] = \det\begin{pmatrix} D_{ee} & 0\\ 0 & D_{ee} \end{pmatrix} \det\begin{pmatrix} 1 & D_{ee}^{-1}D_{eo}\\ D^{-1}D_{ee} & 1 \end{pmatrix} = \det[D_{ee}]\det[D_{oo}]\det[1 - D_{ee}^{-1}D_{eo}D_{oo}^{-1}D_{oe}]$ $= \det[D_{ee}] \det[D_{oo}] \det[\hat{D}_{ee}]$ \hat{D}_{ee} : Schur complement of D UV mode IR mode ILU preconditing [M. Peardon, hep-lat/0011080] Point / stripe blocking for MG solver, Overlap kernel

[A. Boriçi, hep-lat/0704.2341; LAT2007]₁₆

3. Algorithmic developments...(cont'd)

(1) Transform/split det[D] using preconditioner (Action Prec.)

(C) *n*-th root trick and Rational approximation RHMC

[M.Clark, Ph. de Forcrand, A. Kennedy, LAT2005; M. Clark, A.Kennedy, PRL98(2007), PRD75(2007)]

n

$$det[D^{\dagger}D] = det[M] = \left(det[M^{1/n}]\right)^{n} \qquad Action = \sum_{j=1}^{n} \phi_{j}^{\dagger}M^{-1/n}\phi_{j}$$

$$M^{1/n} \qquad \text{UV mode suppressed}$$

$$Action = \phi^{\dagger}M^{-1/n}\phi = \sum_{j=1}^{p} \phi^{\dagger}\frac{\alpha_{j}}{M+\beta_{j}}\phi$$

$$= \sum_{j=\text{UVpole}} \phi^{\dagger}\frac{\alpha_{j}}{M+\beta_{j}}\phi + \sum_{j=\text{IRpole}} \phi^{\dagger}\frac{\alpha_{j}}{M+\beta_{j}}\phi$$

$$M^{-1/n} = \sum_{j=1}^{p} \frac{\alpha_{j}}{M+\beta_{j}}$$

$$M^{-1/n} = \sum_{j=1}^{p} \frac{\alpha_{j}}{M+\beta_{j}}$$

UV mode:large β shift. Large MD Force, small cost IR mode: small β shift. Small MD Force, expensive cost

Distinctive feature: Implicit scale splitting by Rational Approx. 17

(2) MD integrator improvements

Omelyan integrator

[Takaishi & de Forcrand,PRE73(2006); Omelyan, Mryglod & Folk,CPC151(2003)]

 $Q(\delta t)(p,q) = (p,q + \delta t \cdot p)$: evolve q.

 $P(\delta t)(p,q) = (p + \delta t \cdot F, q)$: evolve p.

$$\exp(t\hat{L}_{H}) \approx \left[Q\left(\frac{\lambda t}{n}\right) P\left(\frac{t}{2n}\right) Q\left(\frac{(1-2\lambda)t}{n}\right) P\left(\frac{t}{2n}\right) Q\left(\frac{\lambda t}{n}\right) \right]^{n} = \exp(t\hat{L}_{H'})$$

Shadow Hamiltonian H' (via Baker-Campbell-Hausdorff formula)

$$H' = H + (\alpha \{T, \{T, V\}\} + \beta \{V, \{T, V\}\})\delta t^{2} + O(\delta t^{4})$$
$$\alpha = \frac{6\lambda^{2} - 6\lambda + 1}{12}, \ \beta = \frac{1 - 6\lambda}{24}$$

• Omelyan et al. minimize $\alpha^2 + \beta^2$

$$\lambda_{omelyan} = 0.19318332.....$$

Omelyan integrator / 2nd order Minimum Norm integrator (2MN)

12

50% improvement is observed for QCD (Takaishi & de Forcrand)

24

3. Algorithmic developments...(cont'd)

(2) MD integrator improvements

Extension to Multiple time step integrator for Omelyan

Nesting the Kernel (QPQPQ), K-time scale (depth K)

$$\exp(t\hat{L}_{H}) \approx U_{K-1}(t, (n_0, n_1, \dots, n_{K-1}))$$

 $U_{j}(t,(n_{0},...,n_{j})) = \left[U_{j-1}\left(\frac{\lambda_{j}t}{n_{j}},(n_{0},...,n_{j-1})\right)P_{j}\left(\frac{t}{2n_{j}}\right)U_{j-1}\left(\frac{(1-2\lambda_{j})t}{n_{j}},(n_{0},...,n_{j-1})\right)P_{j}\left(\frac{t}{2n_{j}}\right)U_{j-1}\left(\frac{\lambda_{j}t}{n_{j}},(n_{0},...,n_{j-1})\right)\right]^{n_{j}}$ Recursively defined.

RBC+UKQCD, BMW, QCDSF, ...

 λ_i : tunable parameters

Optimize / Customize your MD integrator

- Shadow Hamiltonian contains errors expressed with Poisson brackets.
- Offline measurement of Poisson brackets;
 Takaishi & de Forcrand, PRE73 (2006);

 Offline measurement of Poisson brackets;
 Clark & Kennedy, LAT2007;

 exp. val.
 <{A,{B,{....}}}</td>
- Minimize the errors by tuning integration parameter, λ, number of time scale, number of pseudo-fermions, ... etc.

- 3. Algorithmic developments...(cont'd)
 - Combination of the UV/IR mode separation and the Multiple time scale MD integrator is now common technique.
 - There still remains the room to improve
 - ○UV/IR separation
 - Blocking, Rational Approx, Preconditiner
 - Low / IR mode : reweighting / Noisy Metropolis
 - OMD integrator
 - Omelyan + Multiple time scale
 - Custom made MD integrator

- 3. Algorithmic developments...(cont'd)
 - Solver Improvements

(1) Mixed Precision / inner-outer solver

- Single precision : effectively doubles memory band width, data cache size, register size.
- Efficiency: S.P. > D.P. Case, mixed prec. is important.
- Intel 64/AMD 64; Single prec. > Double prec.
- Cell PS3/GPGPU; Single >> Double.
- (2) Deflation Technique
 - Remove / suppress small eigenvalues. Better solover behavior
 - Luscher's local coherency for low modes. RG blocking like deflation.
- (3) Multi Grid solver

Adoptive Multi Grid (RG blocking) solver/preconditioner

(1) Mixed precision / Inner-Outer solver Flexible Preconditioner

• Any iterative solver for Ax=b. (short recurrence solver)

[*r* and *x* satisfy r = b - Ax.] [given a scalar " α " and a pre-search vector "*p*".] q = Ap $r = r - \alpha q$ $x = x + \alpha p$ [new *r* and *x* still satisfy r = b - Ax.]

- Accumulated r and x should satisfy r=b-Ax at each update point.
- To make flexible precondition, modify the update lines as

(1) Mixed precision / Inner-Outer solver

• *Right* preconditioning ; AMy = b; x = My.

[*r* and *x* satisfy r = b - Ax.]

[given a scalar " α " and a pre-search vector "p".]

v = Mp q = Av (= AMp: search vector for AMy = b) $r = r - \alpha q$ $x = x + \alpha v$ [new r and x still satisfy r = b - Ax.]

- Search vector is computed for AMy=b.
- The solution-residual relation is kept for r=b-Ax locally.
- This enables us to change M from iteration to iteration (*Flexible* preconditioner).
- Put inner solver for $M pprox A^{-1}$
- M can be single precision. r=b-Ax is kept in double precision. 23

(1) Mixed precision / Inner-Outer solver

- OCG, BiCGStab, CGS,, can be flexible.
 - The most simple case : Richardson / Iterative refinement. [Numerical Recipes]
 - BMW collab. uses D.P. Richardson for outer-solver + S.P. CG for inner-solver

[BMW collab., Dürr et al.,hep-lat/0802.2706]

- PACS-CS: uses D.P. BiCGStab+ S.P. BiCGStab
- For Arnoldi type solver [GMRES,GCR…]
 - Longer reccurence relation
 - Keep a series of intermediate vectors (like v in prev. page.)
 - Then FGMRES, GCR(Lüscher) can be flexible.

○ By tuning solver parameters

- Most Time is spent in (inner) single precision arithmetic.
- If the single precision kernel has much better performance than that with double precision kernel.
- Best performance is obtained with mixed precision solver.
- Promising for GPGPU / CELL computing!!

This is already common to Overlap fermions? Low prec. sign func. (inner) + High prec. sign func. (outer)

(2) Deflation technique

 Critical Slowing down of Solver iteration is caused by small / near zero eigenvalues.

 By subtracting such modes from the matrix spectra, we can recover from the slowding down.

 Deflation technique remove/suppress the near zero eigenspace of *D*.

> This is already common to Overlap fermions (sign functoin)

3. Algorithmic developments...(cont'd)

(2) Deflation technique (To Solve:

 Matrix A has *p*-dimensional subspace with small eigenvalues.

Let c and u spans the subspace.

$$U_{p} = (u_{1}, u_{2}, \dots, u_{p})$$
$$C_{p} = (c_{1}, c_{2}, \dots, c_{p})$$
$$AU_{p} = C_{p}$$
$$C_{p}^{\dagger}C_{p} = I_{p}$$

Suppose the projection operator:

$$P = I_p - C_p C_p^{\dagger}$$
$$Q = I_p - U_p C_p^{\dagger} A$$
$$PA = AQ$$

 Then consider the following preconditiond problem.

$$(PA)y = Pb \qquad \cdots (2)$$

 $Ax = b \qquad \cdots (1)$

 The soluton x of Eq.(1) can be written with y of Eq.(2) as

$$x = Qy + U_p C_p^{\dagger} b$$

$$\therefore Ax = AQy + AU_p C_p^{\dagger} b$$
$$= PAy + C_p C_p^{\dagger} b$$
$$= Pb + C_p C_p^{\dagger} b = b$$

- Solving Eq.(2) is easier than solving Eq.(1), because the coeffcient matrix of Eq.(2) PA does not contains small eigenvalues.
- If The cost to obtain C and U is small, deflation improves solver perfomance.
- How to construct the subspace " C_{ρ} "?

3. Algorithmic developments...(cont'd)

(2) Deflation technique (cont'd)

Many works by Many works by [Luescher, JHEP07(2007),hep-lat/0710.5417; A.Stathopoulos, K.Orginos, hep-lat/0707.0131; W.Wilcox, PoS(LATTICE2007),hep-lat/0710.1813; A.Abdel-Rehim,R.B.Morgan,W.Wilcox,PoS(LATTICE2007); R.B.Morgan,W.Wilcox,math-ph/0707.0505,math-ph/0405053; M.L.Parks, E.De Sturler et al, SIAM J. on Sci.Comp. 28(2006)1651 LATTICE2008: Poster by Abdel-Rehim, Talk by Wilcox]

To avoid exact eigen pairs computation

(a) Overlap eigen mode computation and D^-1 computation.

- GMRES-DR,GMRES-E..:Wilcox, Morgan & Abdel-Rehim
- GCRO-DR: Parks & Sturler

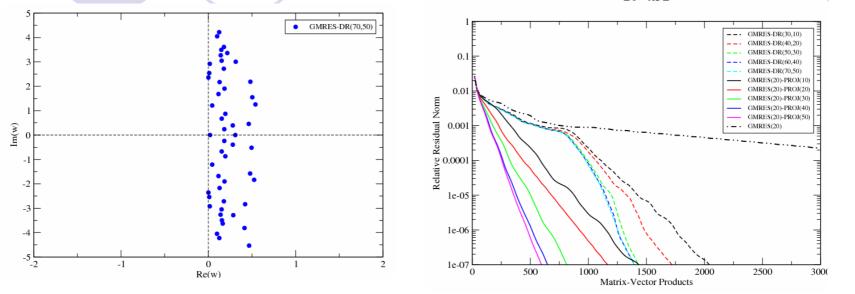
These algorithms can solve *Dx*=*b* and eigen pairs simultaneously.

(b) Make use of Local coherency property of low modes.

Luscher's Domain decomposed subspace blocking with local coherency.

3. Algorithmic developments...(cont'd)

(a) Overlap eigen mode computation and D^-1 computation. Very effective for few Near zero modes / negative eigen modes case. $20^{^{3}}x^{32}$



Near zero modes case

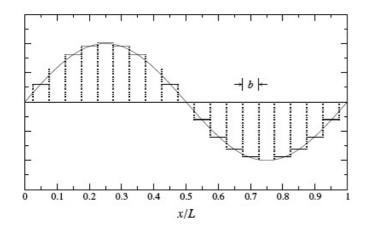
[Wilcox, LAT2007]

- First equation or few equations are solved with GMRES-DR. Once the subspace converged, change solver with GMRES-proj, or Deflated solver.
- Normal GMRES stagnates [dot-dot-dashed line]
- Solver with Deflation/Projection converges. [other lines]
- Critical slowing down is avoided.

[PACS-CS collab. uses GCRO-DR for inner solver] ²⁸

3. Algorithmic developments...(cont'd)

(b) Make use of Local coherency property of low modes. Low modes can be well approximated by few blocked basis vectors [Local coherency]. [Lüscher, JHEP07(2007)081]

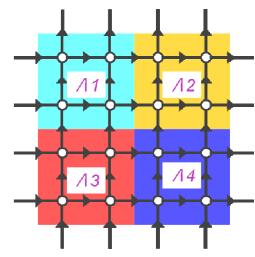


a low mede vector :
$$\psi(x) \approx \sum_{\Lambda}^{blocks} \sum_{j=1}^{N} c_{(j,\Lambda)} \phi_{j}^{\Lambda}(x)$$

$$C = \left\{ \phi_{j}^{\Lambda}(x) : j = 1, \dots, N, \Lambda = \text{all domain blocks} \right\}$$

$$\phi_{j}^{\Lambda}(x) = \begin{cases} \neq 0 \quad (x \in \Lambda) \\ = 0 \quad (x \notin \Lambda), \end{cases} \quad (\phi_{i}^{\Lambda})^{\dagger} \cdot \phi_{j}^{\Lambda'} = \delta_{ij} \delta_{\Lambda\Lambda'}$$

- ϕ is constructed after few smoothing processes via inverse iteration on *N*-random vectors.
- Then blocked and orthogonalized.
 The subspace dimension is effectively enlarged: N x [#of Lattice blocks]
- $C = \{\phi\}$ spans the deflation subspace.
- Suitable for Domain-Decomposition and Memory efficient.



(b) Make use of Local coherency property of low modes.

 Using the Low mode rich subspace C, the deflation projector is constructed as

$$P = 1 - DCB^{-1}C^{\dagger}, \quad Q = 1 - CB^{-1}C^{\dagger}D, \quad B = C^{\dagger}DC,$$

$$PD = DQ, P^2 = P, Q^2 = Q,$$

This contains *B* which is the projection of *D* in to the subspace *C*.

- For Wilson-Dirac operator, the small Wilson-Dirac operator *B* becoms $B(i,\Lambda; j,\Lambda') \equiv \left\langle \phi_i^{\Lambda} \mid D \mid \phi_j^{\Lambda'} \right\rangle$ $= B(i, j,\Lambda) \delta_{\Lambda,\Lambda'} + \sum_{\mu=1}^{4} \left[B(i, j,\Lambda,\mu) \delta_{\Lambda+\hat{\mu},\Lambda'} + B(i, j,\Lambda,\mu) \delta_{\Lambda-\hat{\mu},\Lambda'} \right]$
- Similar to RG blocked W.D.operator. Still has nearest neighbor interaction.
- Using this projection, critical slowing down is avoided.

[Lüscher, JHEP07(2007)081]

3. Algorithmic developments...(cont'd)

(3) MultiGrid Solver

Talk by M. Clark @ this conference. [Brannick,Brower,Clark,Osborn,Rebbi, PRL100(2008);LAT07]

- MultiGrid solver also removes critical slowing down.
- Choice of subspace basis is important. (Prolongator)
- Similar to Luscher's deflation. Low mode enhancement is important.

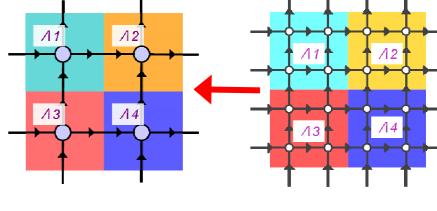
$$v_{l}(x): N \text{ random vector}$$

$$w_{l} = (D)^{-k} v_{l}: \text{ low mode enhanced}$$

$$\text{then } w_{l} \text{ is blocked as}$$

$$C = \left\{ \phi_{j}^{\Lambda}(x): j = 1, \dots, N, \Lambda = \text{all domain blocks} \right\}$$

$$\phi_{j}^{\Lambda}(x) = \begin{cases} \neq 0 \quad (x \in \Lambda) \\ = 0 \quad (x \notin \Lambda), \end{cases} \quad (\phi_{i}^{\Lambda})^{\dagger} \cdot \phi_{j}^{\Lambda'} = \delta_{ij} \delta_{\Lambda\Lambda'}$$



 \odot To solve Dx = b , use the preconditione defined by

$$P \equiv CB^{-1}C^{\dagger} \text{ with } \begin{array}{c} B \equiv C^{\dagger}DC, \text{ or} \\ B(i,\Lambda;j,\Lambda') \equiv \left\langle \phi_i^{\Lambda} \mid D \mid \phi_j^{\Lambda'} \right\rangle \end{array}$$

3. Algorithmic developments...(cont'd)

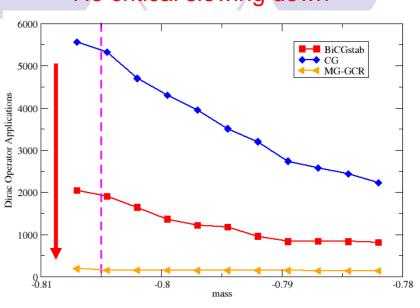
- (3) MultiGrid Solver (cont'd)
- Then Solve

PDx = Pb

- P is the approximation of D^-1 in the subspace C.
- P contains B^-1. to solve this next blocking is applicable.
- Recursively applying this blocking.
 - ⇒ MultiGrid. V cycle
- Similar to Luscher's deflation subspace blocking. Low mode enhancement is important.

QCD 16^3x32 Wilson Case Talk by M. Clark @ this conference. [Brannick,Brower,Clark,Osborn,Rebbi, PRL100(2008);LAT07]

Another RG blocking by A. Borici, hep-lat/0704.2341; LAT2007.



Mass

No critical slowing down

Solver Works in this conference: •J.Bloch (for Overlap fermion) [Mon. Chesapeake C], •J.Osbon (Initial guess for multi-shift solv.)[Mon. Chesapeake C] •W.Wilcox (Deflation/Lanczos/multiple)[Mon. Chesapeake C] •A.Abdel-Rehim (Seed method/multiple) [Poster S.A]

- Mixed precision solver effectively enhances the solver performance.
 - application to GPGPU/CELL?
- Deflation and MultiGrid blocking with low mode-rich basis vector removes Critical slowing down.

Algorithm works in this conference:
[July 15, Tue. Chesapeake B]
A. Bazavov (for HISQ action dynamical sim.)
R.C. Brower (Mobius Algorithm for DW/GapDW fermion.)
M. Clark (Remove Critical Slowing down)
T. Kruth (Dynamically Smeared Fermions)

[July 16, Wed. Chesapeake B]

- •O. Witzel (Polynomial HMC)
- R. Renfrew (Reduce Ch.Sym.breaking for DW)
- F. Palombi (Reweighting for Low mode Quark determinant)
- •W. Cherrington (Dual Lattice Algorithm)
- J. Mucci (SiCortex Machines)

[July 15, Tue. Poster session]

- A. Pochinsky (Efficient QCD code made simpler: qa0)
- •L. Piccoli (Tracking QCD workflows)
- •G. von Hippel (Petrurbative imp. with HISQ fermions)

4. Outlook: Physics at 1PFlops Dynamical QCD simulation at 1 PFlops

• Physical quark masses (*Mud* < 10 MeV, L=3fm,a=0.1fm)

Cost O(10) Tflops Years Wilson/KS type

[ALPHA,BMW,CERN,ETM,JLAB,PACS-CS,QCDSF,MILC,..]

O(100) Tflops Years? Overlap/DW type

[UKQCD/RBC,JLQCD/TWQCD,SESAM/QCDSF,...]

 $(am_c)^2 \approx (1.5/6)^2 = 0.06$ (6% error)

 \bigcirc Finer lattice spacing (1/a > 6 GeV?, L=2fm, 64^3x128 lattice)

Charm quarks

$$\begin{split} \Lambda_{_{QCD}} &\approx 0.3 \text{GeV} < m_{_{charm}} \approx 1.5 \text{GeV} < 1/a \approx 6 \text{GeV} \\ 1/\Lambda_{_{QCD}} &\approx 0.6 \text{fm} > 1/m_{_{charm}} \approx 0.13 \text{fm} > a \approx 0.03 \text{fm} \end{split}$$

Continuum limit

○ Larger lattice volume (L > 6 fm?, 1/a=2GeV, $64^{3}x128$ lattice)

Multi hadron system

Multi scale physics

$$\begin{split} m_{\pi} &\approx 0.1 {\rm GeV} < \quad \Lambda_{_{QCD}} \approx 0.3 {\rm GeV} < 1/a \approx 2 {\rm GeV} \\ 1/m_{\pi} &\approx 2 {\rm fm} \quad > \ 1/\Lambda_{_{QCD}} \approx 0.6 {\rm fm} \quad > \ a \approx 0.1 {\rm fm} \end{split}$$

4. Outlook: Physics at 1PFlops

Dynamical QCD simulation at 1 PFlops

Empirical cost formula

$$\operatorname{Cost}[\operatorname{TFlopsYears}] = C \left[\frac{\#\operatorname{Conf}}{100} \right] \cdot \left[\frac{20 \operatorname{MeV}}{\overline{m}_q} \right]^3 \qquad \cdot \left[\frac{L}{3 \operatorname{fm}} \right]^5 \cdot \left[\frac{0.1 \operatorname{fm}}{a} \right]^7 \qquad [Ukawa, Lat2001 @ Berlin]$$
$$\operatorname{Cost}[\operatorname{TFlopsYears}] = K \left[\frac{\#\operatorname{Conf}}{100} \right] \cdot \left[\frac{20 \operatorname{MeV}}{\overline{m}_q} \right]^1 \qquad \cdot \left[\frac{L}{3 \operatorname{fm}} \right]^5 \cdot \left[\frac{0.1 \operatorname{fm}}{a} \right]^6$$

Now O(10)TFlopsYears for

[DDHMC: Del Debbio et al..JHEP0702(2007)056]

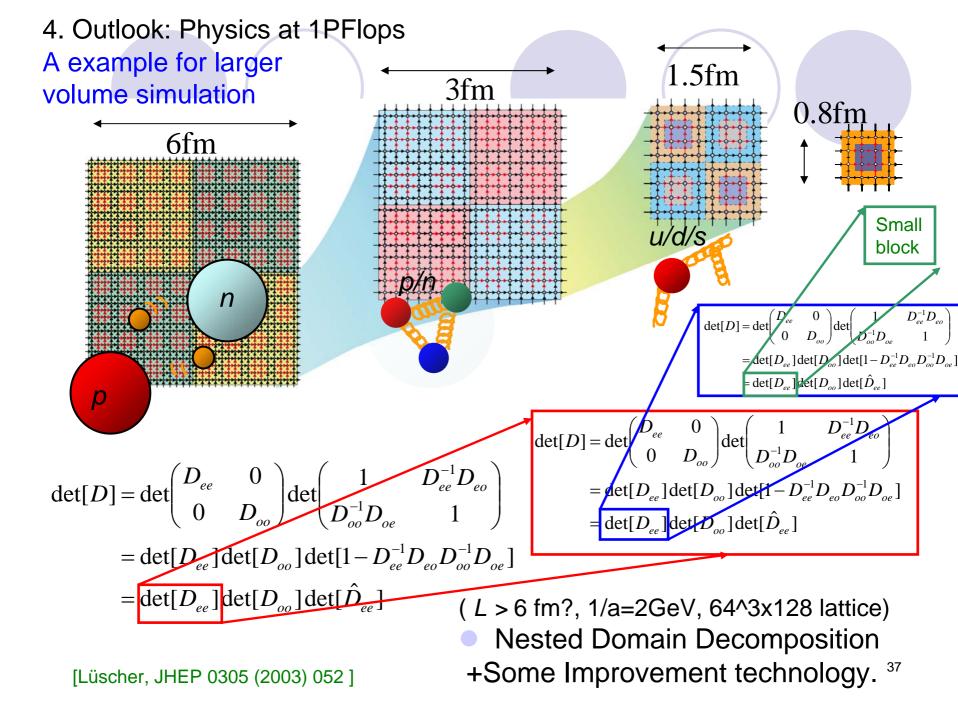
$$M_{\pi}/M_{\rho} \approx 0.2$$
 at $a \approx 0.1 \text{fm}, L \approx 3 \text{fm}$

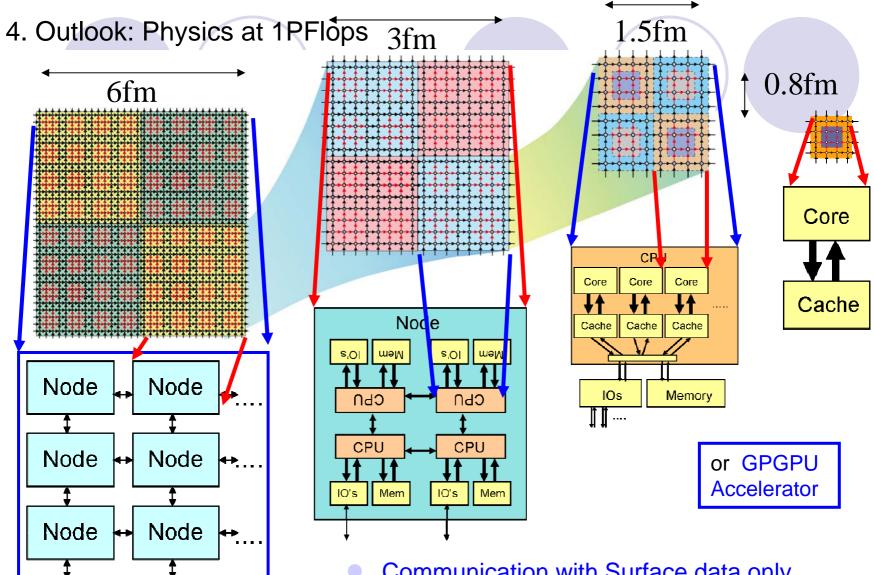
[Talk by Kuramashi, PACS-CS]

- Finer lattice spacing (1/a > 6 GeV?, L=2fm, 64^3x128 lattice)
 - Charm quark on fine lattcie requires a=0.03 fm lattice. The Cost is (2/3)^5*3^6 =96 larger. ⇒ O(1) Pflops Years is required.
 Still difficult problem? ⇒ 10PFlops probrem.
- Larger lattice volume (L > 6 fm?, 1/a=2GeV, $64^{3}x128$ lattice)
 - Multi hadron system by doubling the lattice extent.

The Cost is $2^5 = 32$ larger. $\Rightarrow O(300)$ Tflops Years is required. 30% sustained speed with 1 PFlops peak speed machine can handle this problem.

[Wilcon/KS typo]



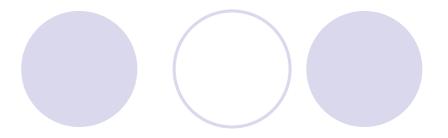


- Communication with Surface data only.
- Bandwidth can be properly treated by this blocking.
- But Latency is limited by speed of light. 38

5. Summary

- Machine trends
 - Multi core architecture is the trend.
 - GPGPU has better cost performance, but actual application for LQCD is now beginning. Large scale simulation is still missing.
 - O CELL becomes common archtecture for HPC?
- Algorithm
 - UV/IR separation + multiple time step MD is common.
 - Deflation and MG remove critical slowing down.
- Physics at 1 PFlops
 - Large volume simulation for multi hadron system can be a target. [Multi scale physics]
 - To tread Multi scale physics, the structure of machine architecture should be taken account.

Backup slides



4. Outlook: Physics at 1PFlops

- Wilson/KS type fermion can handle multi-hadron system with 1Pflops machine in principle.
- Whole System performance analysis that has been done, for ex.
 QCDOC, CP-PACS, APE...., is *again* required.
- Domain-Wall / Overlap fermion : Are there this kind of decomposition ?
- D.W. / 5D-rep. Overlap can use geometric preconditoner.
- 4D-Overlap requires special kernel for geometric decomposition?
 Dirichlet boundary condition for OV op. [Luscher, "Shrodinger Functional with exact Chiral

symmetory", JHEP 0605 (2006) 042]

Enormous works for Dynamical Overlap/DW fermions

[Many people ,RBC,QCDSF,SESAM,JLQCD,.....]

QCD Software / infrastructure works

[MILC code; ILDG; B.Joo,USQCD; A.Borici,QCDLAB; ...]

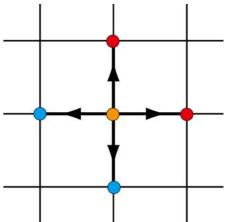
That's all Thank you!

- 2. Machine Trends (cont'd)
- For QCD (dynamical)
 - Hybrid Monte Carlo (HMC)
 - Dynamical Quark part requires huge amount of hopping matrix multiplication.

$$M(n,m) = \sum_{\mu=1}^{4} \left[\left(1 - \gamma_{\mu} \right) U_{\mu}(n) \delta_{n+\hat{\mu},m} + \left(1 + \gamma_{\mu} \right) U_{\mu}^{\dagger}(m) \delta_{n-\hat{\mu},m} \right]$$

○ This computation requires

~ 3 Byte/Flop for a site...

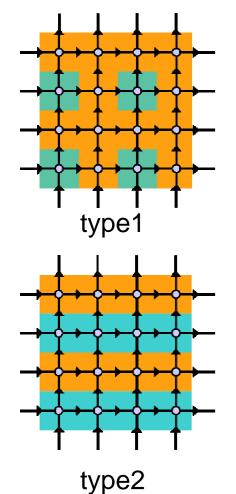


Register, Cache, are memory blocking are required at each layer.

3. Algorithmic developments...(cont'd)

(1) Transform/split det[D] using preconditioner (Action Prec.)(b') Point / stripe (RG) blocking for MG solver, Overlap kernel

[A. Borici, hep-lat/0704.2341; LAT2007]



Change Site Ordering

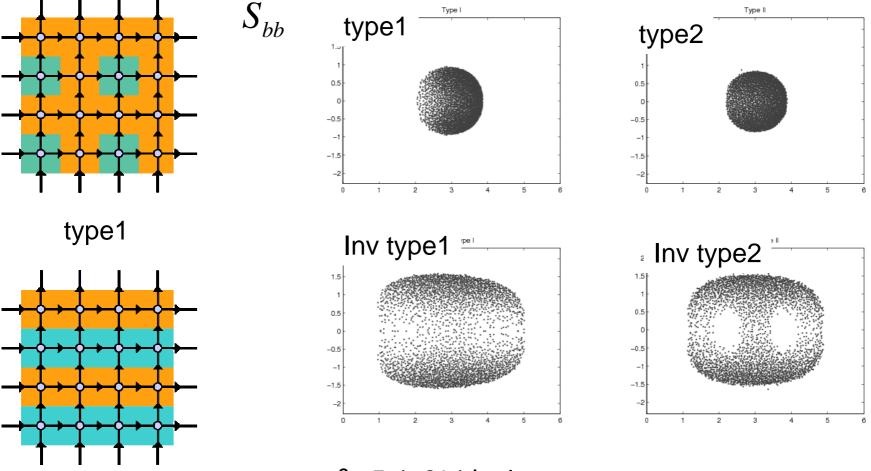
$$det[D] = det \begin{pmatrix} D_{bb} & D_{br} \\ D_{rb} & D_{rr} \end{pmatrix}$$

= $det[D_{rr}]det[D_{bb} - D_{br}D_{rr}^{-1}D_{rb}]$
= $det[D_{rr}]det[S_{bb}]$
UV mode IR mode
 S_{bb} : Schur complement of D

3. Algorithmic developments...(cont'd)

(b') Point / stripe (RG) blocking for MG solver, Overlap kernel

[A.Borici, hep-lat/0704.2341; LAT2007]



type2

3. Algorithmic developments...(cont'd)

(c) *n*-th root trick and Rational approximation RHMC Rational approximation

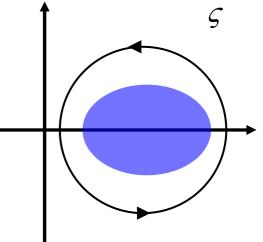
- C.f. Multi boson algorithm
 - $1/x \approx \sum c_i x^j$ Hermitian Polynomial approx. (Luscher '93)
 - Non-Hermitian Polynomial approx. (Borrelli, de Forcrand, Galli ⁽⁹⁶⁾

For RHMC algorithm, similar variant is possible.

Hermitian Rational approx. VS Non-Hermitian Rational approx.

$$D^{-1/n} = \frac{1}{2\pi i} \oint \varsigma^{-1/n} (\varsigma - D)^{-1} d\varsigma$$
$$\approx \sum_{j} (\varsigma_{j})^{-1/n} (\varsigma_{j} - D)^{-1} d\varsigma_{j}$$

D: Non - Hermitian Wilson Dirac op.



3. Algorithmic developments...(cont'd)

(2) MD integrator improvements

- Omelyan integrator
- Simple leapfrog

[Takaishi & de Forcrand, PRE73(2006); Omelyan, Mryglod & Folk, CPC151(2003)]

H(p,q) = T(p) + V(q), : Hamiltonian

$$\begin{pmatrix} p(t) \\ q(t) \end{pmatrix} = \exp(t\hat{L}_{H}) \begin{pmatrix} p(0) \\ q(0) \end{pmatrix}$$
$$\exp(t\hat{L}_{H}) \approx \left[Q\left(\frac{t}{2n}\right) P\left(\frac{t}{n}\right) Q\left(\frac{t}{2n}\right) \right]^{n}$$

 $\hat{L}_{H}X \equiv \{X, H\} = \frac{\partial X}{\partial q} \frac{\partial H}{\partial p} - \frac{\partial H}{\partial q} \frac{\partial X}{\partial p}$ $Q(t) \equiv \exp(t\hat{L}_{T}) : \text{evolve } q.$ $P(t) \equiv \exp(t\hat{L}_{V}) : \text{evolve } p.$ Leapfrog integrator

This operator does not conserve H, but conserves Shadow Hamiltonian H'. $\left[Q\left(\frac{t}{2n}\right)P\left(\frac{t}{n}\right)Q\left(\frac{t}{2n}\right)\right]^{n} = \exp(t\hat{L}_{H'})$ Exact rel.

Shadow Hamiltonian: $H' = H - \frac{1}{24} (\{T, \{T, V\}\} + 2\{V, \{T, V\}\}) \delta t^2 + O(\delta t^4)$

- 3. Algorithmic developments...(cont'd)
- (2) Deflation technique

LQCD requires thousand of linear equation solution

 $Dx^{(i)} = b^{(i)}, i = 1, 2, 3...$ $D^{(i)}x^{(i)} = b^{(i)}, i = 1, 2, 3..., D^{(i)} \approx D^{(i-1)},$

Multiple right-hand side or chain of linear equations.

Quark propagator

Solver in HMC trajectory

- The reduction of condition number of coefficient matrix D is very effective. Efficient Preconditioning is desired.
- Deflation technique is one of the efficient technique to reduce the condition number.

Deflation remove/suppress small eigenspace of D.

3. Algorithmic developments...(cont'd)

(2) Deflation technique (cont'd)

(a) Overlap eigen mode computation and D^-1 computation.

Use Arnoldi type Solver [Krylov subspace method] for Ax=b.

 $r = b - Ax_{0}: \text{ initial residual}$ $v_{0} = r/|r|,$ $AV_{K} = V_{K+1}\overline{H}_{K}: \text{ Arnoldi factorization (via Gram - Schmidt).}$ $V_{K+1} = (v_{0}, v_{1}, \dots, v_{K+1}), \quad v_{i}^{\dagger}v_{j} = \delta_{ij}, \quad \text{Krylov subspace basis.}$ $\overline{H}_{K}: \text{ upper Hessenberg matrix } (K + 1 \times K)$ $x = x_{0} + V_{K+1}c, \quad \text{where } c^{T} = (c_{0}, c_{1}, \dots, c_{K+1})$ $\text{Minimize: } |r| = |b - Ax| \quad \text{with } c.$

- GMRES(K)
- V_{K+1} and H_K contains the spectrum info. of A.
- At restarting, construct Harmonic-Ritz pairs.

Solve small eigen problem $(K \times K)$

$$(\overline{H}_{K}^{\dagger}\overline{H}_{K}-\mu H_{K}^{\dagger})y=0$$

GMRES-DR/GCRO-DR

Harmonic Ritz pair $(\mu, w = V_K y)$ is approximation for $(A - \mu)w = 0$.

Few $\{w\}$ basis vectors are recycled as deflation subspace for the next iteration.

Reduce eigen mode comp. cost

3. Algorithmic developments...(cont'd)

(b) Make use of Local coherency property of low modes.

Deflation projector contains small linear equation B^-1.

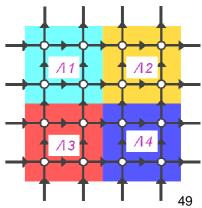
$$P = 1 - DCB^{-1}C^{\dagger}, \quad Q = 1 - CB^{-1}C^{\dagger}D, \quad B = C^{\dagger}DC,$$

$$PD = DQ, P^2 = P, Q^2 = Q,$$

- For Wilson-Dirac operator, the small Wilson-Dirac operator *B* becoms $B(i,\Lambda;j,\Lambda') \equiv (\phi_i^{\Lambda})^{\dagger} \cdot (D\phi_j^{\Lambda'})$ $= B(i,j,\Lambda)\delta_{\Lambda,\Lambda'} + \sum_{\mu=1}^{4} \left[B(i,j,\Lambda,\mu)\delta_{\Lambda+\hat{\mu},\Lambda'} + B(i,j,\Lambda,\mu)\delta_{\Lambda-\hat{\mu},\Lambda'} \right]$
- Similar to RG blocked W.D.operator. Still has nearest neighbor interaction.
- To avoid frequent application of Projection,
 P is applied to SAP preconditioned problem:

$$DM_{SAP} y = b, x = M_{SAP} y,$$

$$\Rightarrow PDM_{SAP} y = Pb, x = QM_{SAP} y + CB^{-1}C^{\dagger}b.$$



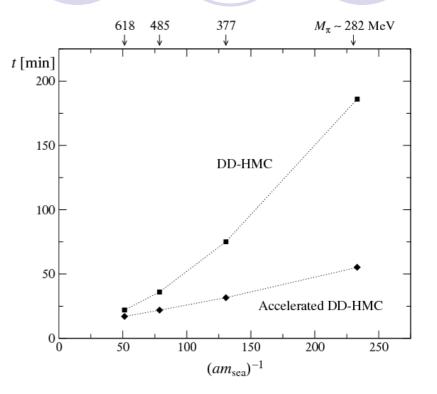
[Lüscher, JHEP07(2007)081]

3. Algorithmic developments...(cont'd)

[Lüscher,hep-lat/0710.5417]

(b) Make use of Local coherency property of low modes.

- Deflation accelerates DDHMC performance
- Factor 2-3 improvement is observed.
- Speedup is significant for smaller quark masses.



Deflation removes critical slowing down.

3. Algorithmic developments...(cont'd)

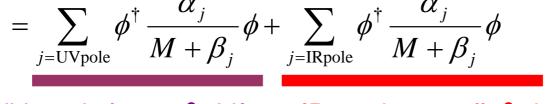
(c) *n*-th root trick and Rational approximation RHMC Further cost reduction using Rational approximation

$$M^{-1/n} = \sum_{j=1}^{p} \frac{\alpha_j}{M + \beta_j}$$

Action =
$$\phi^{\dagger} M^{-1/n} \phi = \sum_{j=1}^{p} \phi^{\dagger} \frac{\alpha_{j}}{M + \beta_{j}} \phi$$

M: Hermitian, spectrum boundary is known. α , β : real parameter Optimal Chebyshev approx.

Partial fraction form



UV mode: large β shift. IR mode: small β shift.

RHMC: RBC+UKQCD, DW Nf=2+1 simulation [hep-lat/0804.0473;PRD76(2007)] Clark and Kennedy KS fermion [hep-lat/0610047;PRD75(2007)] Takaishi and Nakamura, One-flavor Wilson fermion F.T. [LAT2007,hep-lat/0711.3888]

3. Algorithmic developments for dynamical QCD (Wilson type)

Lattice QCD partition function

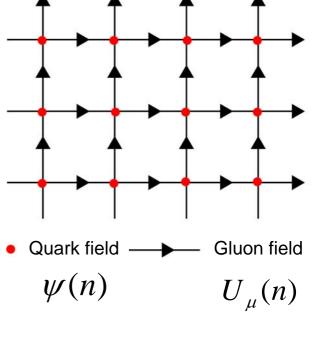
$$Z = \int \prod dU d\psi e^{-S[U,\psi]}$$

 $S[U,\psi] = S_G[U] + S_Q[U,\psi]: \text{Lattice QCD action}$ $S_G[U]: \text{Gluon part} (\rightarrow Tr[F_{\mu\nu}F_{\mu\nu}]/(4g^2))$ $S_Q[U,\psi]: \text{Quark part} (\rightarrow \sum_f \overline{\psi}_f (D+m_f)\psi_f)$

Nf=2+1 partition function
 (ψ integ.out)

$$Z = \int \prod dU \det[D_{ud}]^2 \det[D_s] e^{-S[U]}$$

HMC algorithm to generate $\{U\}$.



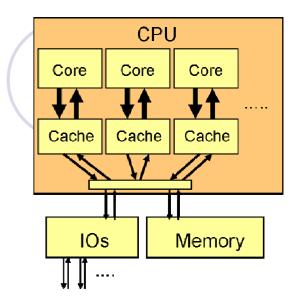
2. Machine Trends (cont'd)

- 1~10 PFlops machine Bottlenecks
 - Memory band width
 - DDR3(1333) 10GB/sec
 - RambusXDR 26GB/sec
 - Byte/Flop < 0.25 (single CPU)</p>
 - GPGPU is more better 100GB/sec

Multi slots/node enhances the node speed. [SMP or NUMA] but...

- IO/ Network band width Depends on the NIC but
 Myrinet 10G 1.25GB/sec
 Myrinet 10G 1.25GB/sec
 - Infiniband DDR 2.0GB/sec
 - Ex. Byte/Flop < 2/48 = 0.04
 - To balance, multi rail (x4 or x8...)

Blocking is required at each level (core/cpu/node) for 1PFlops machine₅₃



 $M(n,m) = \sum_{\mu=1}^{4} \left[(1 - \gamma_{\mu}) U_{\mu}(n) \delta_{n+\hat{\mu},m} + (1 + \gamma_{\mu}) U_{\mu}^{\dagger}(m) \delta_{n-\hat{\mu},m} \right]$ Hopping Mult : $\sim 3Byte/Flop req.$

3. Algorithmic developments...(cont'd)

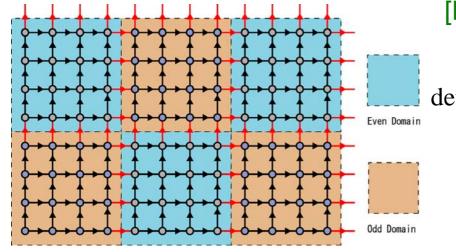
(c) *n*-th root trick and Rational approximation RHMC Further cost reduction using Rational approximation

Action =
$$\phi^{\dagger} M^{-1/n} \phi = \sum_{j=UVpole} \phi^{\dagger} \frac{\alpha_j}{M + \beta_j} \phi + \sum_{j=IRpole} \phi^{\dagger} \frac{\alpha_j}{M + \beta_j} \phi$$

$$\int_{0}^{100} \frac{\phi^{\dagger}}{\alpha_j} \frac{\phi^{\dagger}}{\alpha_j} \frac{\alpha_j}{\alpha_j} \frac{\phi^{\dagger}}{\alpha_j} \frac{\phi^{\dagger}}{$$

IR mode: small β shift Expensive cost, small force => Coarser MD step UV mode:large β shift Cheap cost, large force => finer MD step 54

- 3. Algorithmic developments...(cont'd)
- (1) Transform/split det[D] using preconditioner (Action Prec.)
 - (b) Luscher Domain-Decomposition preconditioned DDHMC



DDHMC simulations: •ALPHA: Von Hippel •CERN: Luscher, Debbio, Giusti, Petronzio •PACS-CS [Lüscher, JHEP 0305 '03, CPC 165 '05]

$$t[D] = det \begin{pmatrix} D_{ee} & 0 \\ 0 & D_{oo} \end{pmatrix} det \begin{pmatrix} 1 & D_{ee}^{-1}D_{eo} \\ D_{oo}^{-1}D_{oe} & 1 \end{pmatrix}$$
$$= det[D_{ee}]det[D_{oo}]det[1 - D_{ee}^{-1}D_{eo}D_{oo}^{-1}D_{oe}]$$
$$= det[D_{ee}]det[D_{oo}]det[\hat{D}_{ee}]$$
$$UV \text{ mode } \qquad \text{IR mode}$$
$$\hat{D}_{ee}: \text{ Schur complement of D}$$

How about another decomposition/blocking?

- ILU preconditing [M. Peardon, hep-lat/0011080]
- Point / stripe blocking for MG solver, Overlap kernel

[A. Boriçi, hep-lat/0704.2341; LAT2007] ⁵⁵

3. Algorithmic developments...(cont'd)

(2) MD integrator improvements Optimize / Customize your MD integrator

Takaishi & de Forcrand, PRE73 (2006); Clark & Kennedy, LAT2007; Poster by Kennedy

- Shadow Hamiltonian contains errors expressed with Poisson brackets.
- Offline measurement of Poisson brackets; exp. val. <{A,{B,{....}}}>
- Minimize the errors by tuning integration parameter, λ , number of time scale, number of pseudo-fermions, ... etc.