Parallel Computing

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Types of "Parallel" Computing

- Nomenclature
 - Parallel vs. concurrentBit-level vs. data level
- Multi-threaded
- Multi-process
- Distributed
- Grid
- SIMD
- GPU/GPGPU
- MIC



SIMD = Single Instruction Multiple Data

- Special registers on CPU where multiple numbers can be packed and operated on simultaneously
- Also known as "vectorization"
 - gcc: "...vectorization is enabled by the flag -ftree-vectorize and by default at -O3"
- CPU vendors have their own implementations and evolutions

(e.g. Intel has ...)

- MMX (1997, Pentium 5) 64bit
- SSE (1999) SSE4(2006) 128 bit
- AVX (2008) 256 bit
- MIC/VPU

512 bit

SIMD = Single Instruction Multiple Data

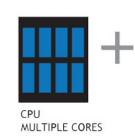
	Constal wastered and CDU sub-ana wastered a wastered
297	<pre>// Multiply a 5x1 matrix by its transpose</pre>
298	inline DMatrix5x5 MultiplyTranspose(const DMatrix5x1 &m1){
299	ALIGNED_16_BLOCK_WITH_PTR(m128d, 5, p)
300	m128d &b1=p[0];
301	m128d &b2=p[1];
302	m128d &b3=p[2];
303	m128d &b4=p[3];
304	m128d &b5=p[4];
305	<pre>b1=_mm_set1_pd(m1(0));</pre>
306	b2=_mm_set1_pd(m1(1));
307	b3=_mm_set1_pd(m1(2));
308	b4=_mm_set1_pd(m1(3));
309	b5=_mm_set1_pd(m1(4));
310	<pre>return DMatrix5x5(_mm_mul_pd(m1.GetV(0),b1),_mm_mul_pd(m1.GetV(0),b2),</pre>
311	_mm_mul_pd(m1.GetV(0),b3),_mm_mul_pd(m1.GetV(0),b4),
312	_mm_mul_pd(m1.GetV(0),b5),
313	_mm_mul_pd(m1.GetV(1),b1),_mm_mul_pd(m1.GetV(1),b2),
314	_mm_mul_pd(m1.GetV(1),b3),_mm_mul_pd(m1.GetV(1),b4),
315	_mm_mul_pd(m1.GetV(1),b5),
316	_mm_mul_pd(m1.GetV(2),b1),_mm_mul_pd(m1.GetV(2),b2),
317	_mm_mul_pd(m1.GetV(2),b3),_mm_mul_pd(m1.GetV(2),b4),
318	_mm_mul_pd(m1.GetV(2),b5));
319	}
	• IVIIC/VPU 512 DIT

GPU = Graphics Processing Unit

- Driven by gaming industry where high frame rates of complex environments was required
- Many cores (few x10¹⁻³) used in lockstep to calculate same algorithm with different inputs
- Programmed via special API
 - CUDA, OpenCL, OpenGL



GPU – Example CUDA code



GPU THOUSANDS OF CORES

Standard C Code

```
void saxpy(int n, float a,
                                float *x, float *y)
{
    for (int i = 0; i < n; ++i)
        y[i] = a*x[i] + y[i];
}
int N = 1<<20;</pre>
```

```
// Perform SAXPY on 1M elements
saxpy(N, 2.0, x, y);
```

C with CUDA extensions

```
__global__
void saxpy(int n, float a,
                             float *x, float *y)
{
    int i = blockIdx.x*blockDim.x + threadIdx.x;
    if (i < n) y[i] = a*x[i] + y[i];
}</pre>
```

int N = 1<<20; cudaMemcpy(x, d_x, N, cudaMemcpyHostToDevice); cudaMemcpy(y, d_y, N, cudaMemcpyHostToDevice);

// Perform SAXPY on 1M elements
saxpy<<<4096,256>>>(N, 2.0, x, y);

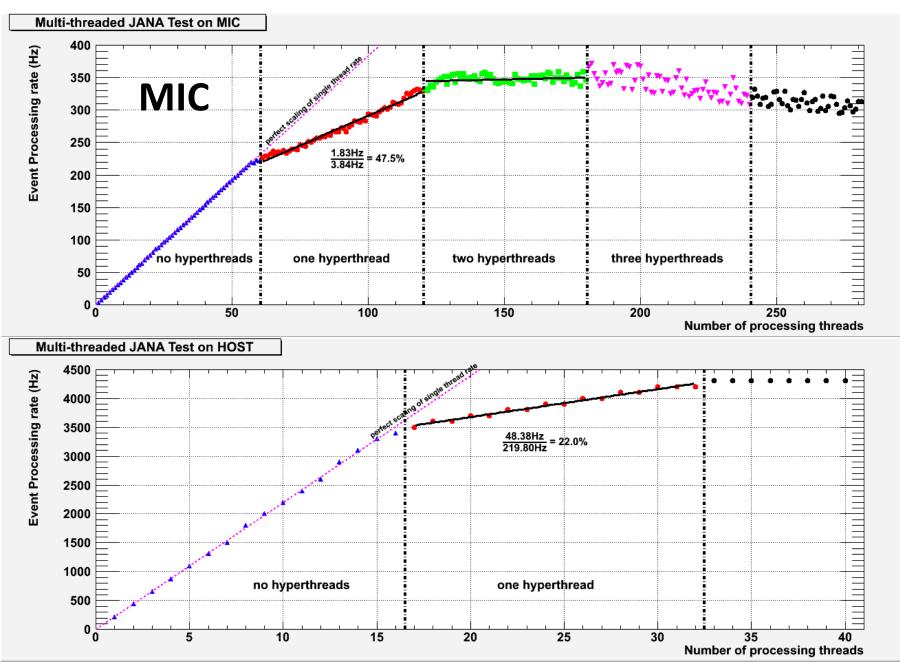
cudaMemcpy(d_y, y, N, cudaMemcpyDeviceToHost);

MIC = Many Integrated Cores

- Xeon Phi = Intel's MIC system
 - 60 cores, 1GHz on a PCIe x16 card
 - 512 bit wide vectors
 - Original project: Larrabee
 - Attempt to make GPU from older x86 design
- Linux variant runs on MIC card independent of host OS
 - MIC system is based on 2.4 Linux kernel
 - File system not automatically shared
 - MIC cards can be configured to mount host's filesystem via NFS
- Must use intel-provided cross-compiler to build executables
 - Could not build sim-recon because ROOT was needed
 - Could not build ROOT because libX11-devel was needed



Performance using JANA TestSpeed plugin



Multi-Threading vs. Multi-Process

<u>Multi-threaded</u>

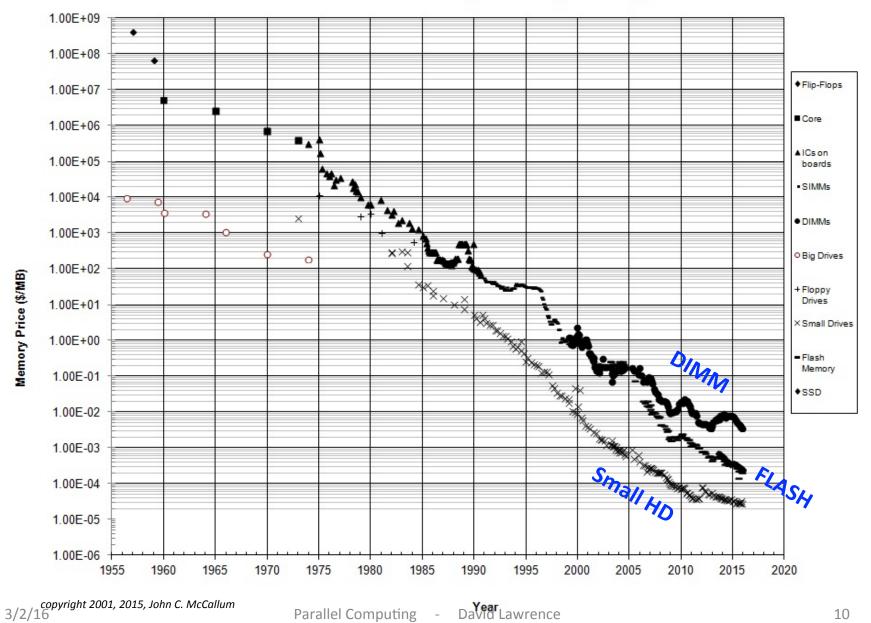
• Uses same memory space

- Uses less RAM
- Minimizes disk head thrashing

Less overhead in process
 management

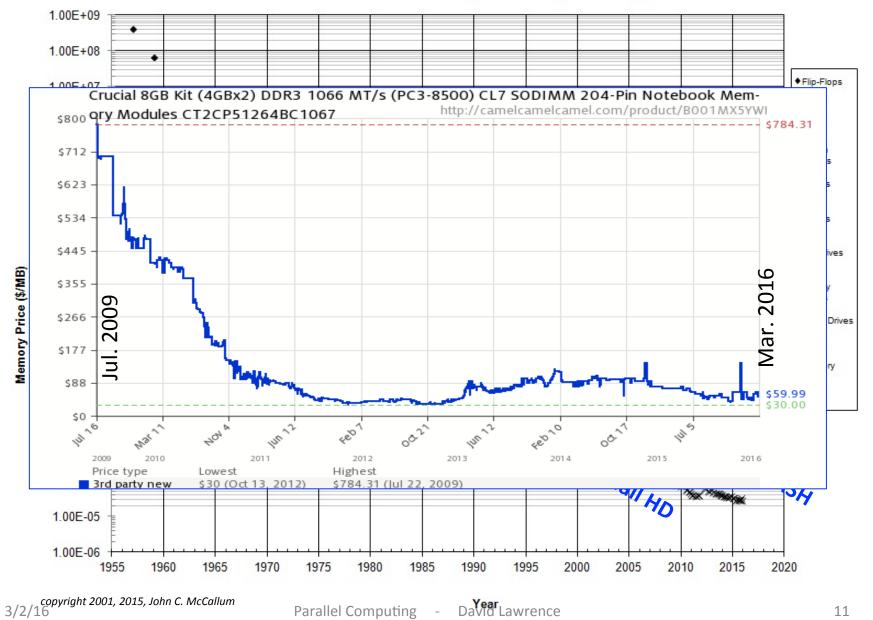
Multi-process

- Uses special shared memory segments or other message passing protocol
- Independent program contexts
- Simpler to debug single thread program crashes
- Less expertise required by novice programmers



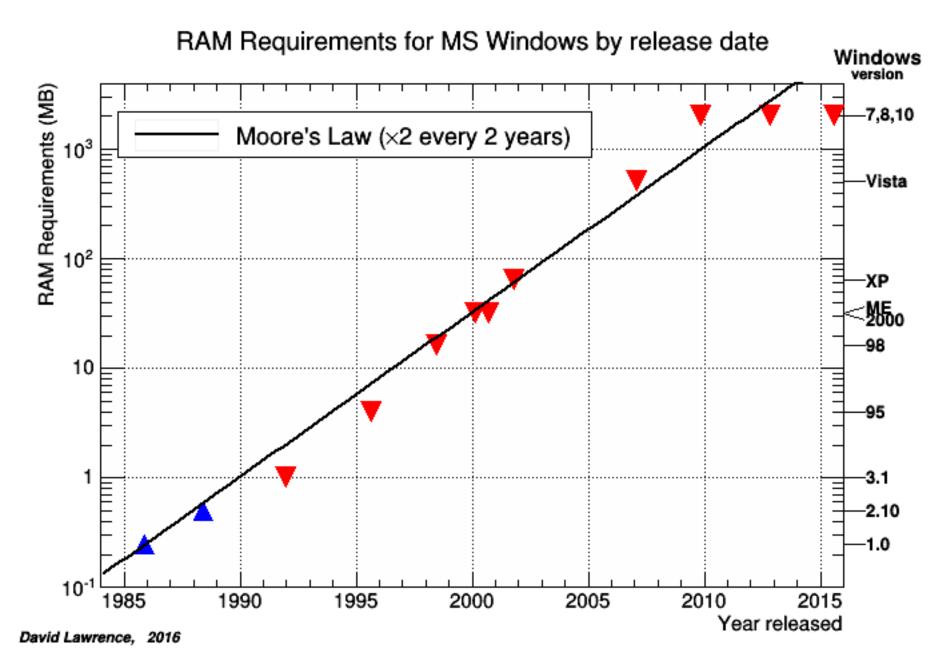
Historical Cost of Computer Memory and Storage

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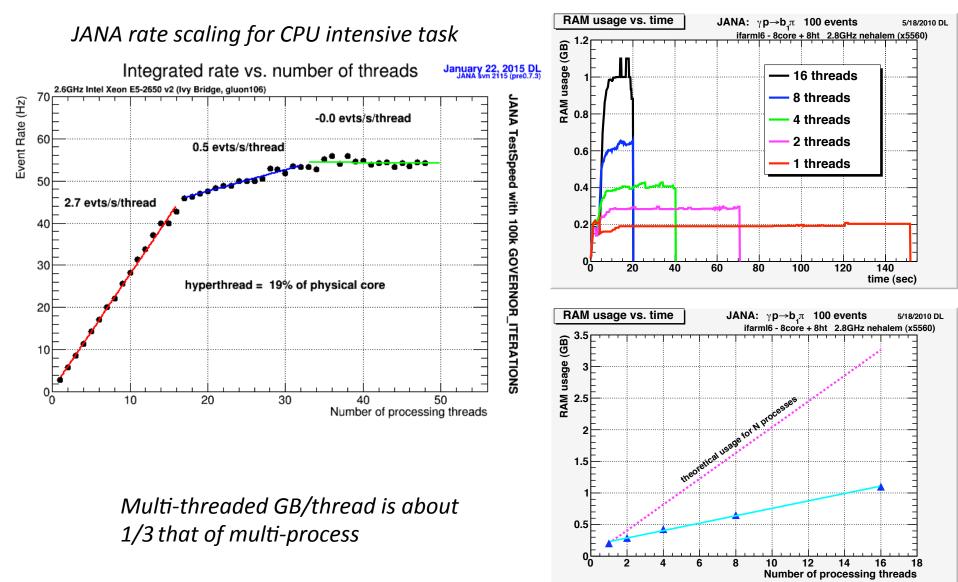


Historical Cost of Computer Memory and Storage

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JANA RAM usage

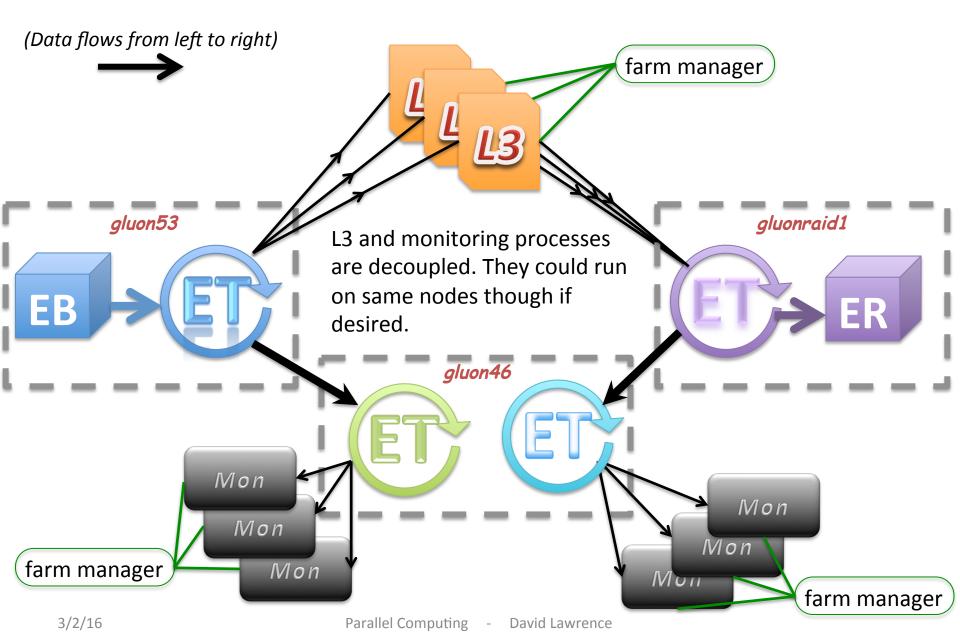


Distributed Computing (Let's just call it "farms")

- large cluster of computers, housed in same location, and connected via fast LAN
- jobs run independently on single node (... or maybe not ...)
- focuses significant compute power to dedicated job
- "clouds" tend to be made up of multiples of these connected via WAN



Hall-D L3 and monitoring architecture



Farms in the Future

Farms will play a role in the future due to power supply and dissipation

(i.e. You can't pack too many teraflops into a small volume without burning everything up!)



Open Science GRID

 US-based federation of compute infrastructures for research and academic communities



Open Science GRID

Characteristics of OSG Jobs

- The application is a Linux application for the x86 or x86_64 architecture.
- The application is single- or multi-threaded but does not require message passing.
- The application has a small runtime between 1 and 24 hours.
- The application can handle being unexpectedly killed and restarted.
- The application is built from software that does not require contact to licensing servers.
- The scientific problem can be described as a workflow consisting of jobs of such kind.

Gulf of

 The scientific problem requires running a very large number of small jobs rather than a few large jobs.

What do big LHC Experiments Do?

- CMS
 - Initially developed reconstruction algorithms as single-threaded
 - Developed framework that can identify which algorithms can be run in parallel giving sub-event level parallelism
- ATLAS
 - Single threaded using GAUDI
 - Developing GAUDIHive for parallelism at event, algorithm, and sub-algorithm levels

What do RHIC Experiments Do?

Who knows?
 (Online documentation is quite old)



What to plan for? (IMHO)

- Parallelism will play a role. Effective systems will likely take advantage of multiple technologies
 - Tough to tell what emerging technologies will stick around. (Be careful, but not overly cautious!)
 - Spend time on the API. It will allow parts of the software to be replaced later without rewriting everything
 - Stay away from requiring commercial products (libraries). Stick with Open Source
 - Commit resources to software R&D. Not just at the beginning, but continuing...

Additional Thoughts (IMLTHO)

- The accessibility of a computing technology should be inversely proportional to the developer base (i.e. if you commit everyone to an obscure or expensive technology, a lot less people we be able to contribute)
- We have an obligation to next generation students/scientists to engage them in technologies that carry them at least a little way into the future as opposed to holding them in our past

Backups



